On the Performance Limits of Joint Source-Channel Coding Using Subband Decomposition

and OFDM Transmission

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Goal: To make a joint source-channel coding (JSSC) system Why? JSSC to make efficient and robust system with low delay How?

- Signal decomposition by subband filtering
- Channel representation by OFDM
- In between: Dimension changing mappings

More or less symmetric system

How does the structure imposed degrade compared to the optimal system?





Transmit vector \mathbf{x} with N components using channel representation \mathbf{y} with K components.

- Bandwidth compression: K < N
- Bandwidth expansion: K > N

Bandwidth compression generally introduces approximation noise

For expansion approximation noise can be avoided

Received and decoded signal: $\hat{\mathbf{x}} = \mathcal{R} \circ (\mathcal{M} \circ \mathbf{x} + \mathbf{n}), \qquad (\mathcal{R} \approx \mathcal{M}^{-1})$



Example where bandwidths of signal and channel are different





Calculating the source *rate-distortion function* at the *channel capacity* renders OPTA (optimal performance theoretically attainable)

Assumptions:

- The signal is Gaussian with a known non-white spectrum
- The channel is additive, Gaussian with a known noise spectrum and attenuation (assumption: constant attenuation)
- The signals and channels are time discrete (no loss in performance due to A/D and D/A processes)



Rate (in bits per second):

$$R(\lambda) = \int_{-W_s}^{W_s} \max\left\{0, \frac{1}{2}\log_2\frac{S_{XX}(F)}{\lambda}\right\} dF$$

Corresponding distortion:

$$D(\lambda) = \int_{-W_s}^{W_s} \min \left\{\lambda, S_{XX}(F)\right\} dF$$

By selecting a value of λ a rate and the corresponding distortion result



The capacity in bits per second is given by

$$C(\theta) = \int_{-W_c}^{W_c} \frac{1}{2} \log_2 \left\{ 1 + \frac{(\theta - S_{NN}(F))^+}{S_{NN}(F)} \right\} dF,$$

where the parameter $\boldsymbol{\theta}$ is found from the power constraint

$$\int_{-W_c}^{W_c} \left(\theta - S_{NN}(F)\right)^+ dF = P$$

The following definition has been used:

$$(x)^{+} = \begin{cases} x \text{ if } x > 0\\ 0 \text{ otherwise} \end{cases}$$



Signal: AR(1) process with $R_{XX}(k) = \sigma_X^2 \rho^{-|k|}$, $\rho = 0.9$ (bandwidth=1)

Channel noise: $S_{NN}(f) = a * (0.1 + f), f \in [0, r]$ (r: channel bandwidth)



White signal and noise:

If the signal has bandwidth B and variance (σ_X^2) higher than the signal noise variance (σ_D^2) , then

$$R = B \log_2 \left(\frac{\sigma_X^2}{\sigma_D^2} \right),$$

and the channel power density (S) is higher than the channel noise density (N), then

$$C = W \log_2\left(1 + \frac{WS}{WN}\right)$$

Equating rate with capacity and solving with respect to the distortion, we obtain

$$\sigma_D^2 = \sigma_X^2 \left(1 + \frac{S}{N} \right)^{-r}, \ (r = W/B)$$



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OPTA does not indicate a structure for making systems with good quality

Shannon's separation theorem states that source coding and channel coding can be done separately without loss of optimality

But:

- Optimal source coding requires infinite dimensional Vector Quantizer
- Optimal channel coding requires e.g. a turbo-like coder with infinite delay and complexity

We will look at alternative structures for getting close to OPTA





- Signal decomposition by transform or filter bank
- Dimension change by nonlinear mapping
- Signal recombination by OFDM



Source decomposition results in several sub-sources with different statistics (mainly different variances)

Each sub-source is to be transmitted on sub-channel where each channel has different statistics

Goal: By constraining the overall bandwidth and the total power, how do we allocate the available resources to the sub-channels?

Result obtained is relevant for a single user, but can be applied to multiuser system if average SNR is the optimization criterion



- Split signal by ideal uniform filter bank into M subbands
 - The split is so fine that the spectrum is constant within each subband
- Map each subband individually to a part of the channel using optimal bandwidth and power in order to fill total channel bandwidth and exactly use allowed power
 - Each part of the channel has constant noise spectrum and attenuation
 - Assume OPTA performance for each mapping



Visby'04

Assume that we know the channel noise level experienced by each signal component

Total channel power:

$$\sum_{m=1}^{M} \sigma_{C_m}^2 = \sum_{m=1}^{M} S_m W_m = \sum_{m=1}^{M} S_m B r_m \le P_{tot}$$

Total bandwidth:

$$\sum_{m=1}^{M} W_m = \sum_{m=1}^{M} Br_m = MBr_{avg} = W_{tot}$$

Optimization using the following object function:

$$O = \sum_{m=1}^{M} \sigma_{D_m}^2 + \lambda_1 \sum_{m=1}^{M} S_m r_m + \lambda_2 \sum_{m=1}^{M} r_m$$

The optimization leads to:

Power density relations:

$$S_j + (S_j + N_j) \log \left(1 + \frac{S_j}{N_j}\right) = S_k + (S_k + N_k) \log \left(1 + \frac{S_k}{N_k}\right)$$

Rates:

$$r_{j} = \frac{c_{a}}{c_{j}} \left[r_{avg} + \frac{1}{2M} \sum_{k=1}^{M} \frac{1}{c_{k}} \log_{2} \left(\frac{\sigma_{X_{k}}^{2}}{\sigma_{X_{j}}^{2}} \frac{N_{j} + S_{j}}{N_{k} + S_{k}} \right) \right].$$

where

$$c_m = \frac{1}{2} \log_2 \left(1 + \frac{S_m}{N_m} \right) \text{ and } \frac{1}{c_a} = \left(\frac{1}{M} \sum_{j=1}^M \frac{1}{c_j} \right)$$

How does water filling relate to this result? Optimal when:

- all the channels have equal bandwidths
- all the sources have equal variances



The present model assumes a constant noise level in each channel

If we have a continuous channel which is split into sub-channels, we do not a priori know what noise densities to use

Strategy:

- Assume a uniform distribution of bandwidths and use the noise density at the midpoint to represent each channel band
- Use the formulas to find the corresponding bandwidths if the assumption were true
- Split the channel into bands according to the obtained bandwidths and find midpoint noise densities
- Calculate new densities and bandwidths
- Continue process until convergence



Signal: AR(1) process with $\rho = 0.9$

Noise: Same or reversed spectrum (Channel capacity equal for both)

Number of channels: 30

Rate (bandwidth) change: r=0.5





Red: Opposite spectra Green: Equal spectra





Relatively small power variations

By inspecting the Lagrange multipliers, the one related to bandwidth dominates

From these observations it is tempting to optimize without the power constraint (but still adjust the power)

Resulting rate formula:

$$r_k = \frac{c_a}{c_k} \left[r_{avg} + \frac{1}{2M} \sum_{j=1}^M \frac{1}{c_j} \log_2 \left(\frac{\sigma_{X_k}^2}{\sigma_{X_j}^2} \frac{c_k}{c_j} \right) \right]$$



1000 channels (close to optimal)



Red squares: Equal signal and noise spectra Blue dots: Noise spectrum turned around



NTNU-

Bandwidths

Opposite spectra, 30 channels





Bandwidths

Equal spectra, 30 channels





Equal spectra, 30 channels







The above theory could be used as an approximation if the OFDM channels are quite narrow, that is each OFDM channel is much smaller than W_m , m = 1, 2, ..., M to get approximately correct bandwidths

On the other hand, the noise must be the same in all the OFDM channels for the theory to be exact

If not, the bandwidths must be quantized while preserving total bandwidth

Only a small set of dimension changing mappings will be available

Practical non-linear mapping will lead to further degradation







Possible channel repr.: Distance from origin used as PAM channel samples:

blue curve: positive amplitudes

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red curve: negative amplitudes
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Loss: approx. 1.5 dB compared to OPTA

Use rates 3, 2, 1, 3/2, 1/2, 1/3, 1/4, 0



Even greater loss using practical mappings



Develop bandwidth allocation procedure when a small set of mappings is available

Look at alternative ways of optimization (Calculus of variations)

Design system: Subband decomposition-mappings-OFDM including adaptive allocation of maps

Channel estimation: Theory requires exact channel knowledge. Use pilots to find approximate channel states and transmit states in feedback channel. Find optimal optimal number of pilots and their power levels

