

Issues in multiuser diversity

- 1. Feedback and real-time traffic
- 2. Traffic prediction and uncertainty
- 3. Multiple antennas

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The benefit of channel variations

Consider a cellular downlink:

- Multiple users sharing the same resources
- Varying channel quality ride the peaks!
 - More variations \Rightarrow more throughput



Multiuser diversity



Limited channel feedback

Consider a downlink using adaptive modulation.

Each timeslot, each mobile feeds back $log_2(M+1)$ bits indicating which rate the channel supports given a desired BER

$$r_u \approx \log_2\left(1 + \frac{\mathrm{SNR}_u}{f(\mathrm{BER}_u)}\right)$$





Quantization for maximum expected throughput

• The expected throughput with rate thresholds q₀...q_M becomes

$$\langle x \rangle = \sum_{m=1}^{M} q_m \left(\prod_{u=1}^{U} \int_0^{q_m+1} p(r_u) dr_u - \prod_{u=1}^{U} \int_0^{q_m} p(r_u) dr_u \right)$$

- *U* is the number of users
- $-p(r_u)$ is the probability that user *u* can receive with rate r_u at a desired BER.
- 1-bit feedback:

$$\langle x \rangle = q \left(1 - \prod_{u=1}^{U} \int_{0}^{q} p(r_u) dr_u \right)$$







As the number of users increases, the throughput becomes extremely sensitive to the choice of *q*.

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Implications of 1-bit channel feedback

- Theoretically, strict multiuser diversity is not badly affected by limited feedback
- but in practice, an <u>extreme sensitivity</u> to correct quantizations leads to drastic performance drops
- Note also that unfairness increases when feedback is reduced



Possible remedies

- Individual thresholds
 - decrease the sensitivity
 - but optimal individual thresholds depend on other users' thresholds
- Avoid using strict multiuser diversity
- Increase the channel feedback

We will combine the two first suggestions and at the same time attain short-term fairness



Diversity-Enhanced Equal Access

- A modified fair multiuser-diversity strategy
 - Round-robin tournament:
 - In each time slot transmit to the user with best channel of the users that have not yet accessed the channel.
 - When all users have obtained access, repeat the tournament.
 - Determine individual 1-bit quantizations locally at the mobile terminal
- Ideal for real-time traffic
- Unsensitive to channel quantization errors



- Note that the scheduler uses a strict multiuser-diversity strategy with a decreasing number of users
- On average, U/2 users compete
 - Performance will on average equal that of strict multiuser diversity with U/2 users
- Ex: Rayleigh fading, 16 users spread out
 - Throughput becomes 268% of that of a round-robin scheme with an optimal common 1-bit quantization







2. Traffic prediction and adaptation

- In theory, better performance can be obtained by scheduling over several time slots
 - Particularly with QoS constraints
- Requires channel prediction over longer periods
 - Should average criterion over pdf for channel to account for higher uncertainty
- Requires traffic prediction
 - "Always data to send" unrealistic assumption



Scheduling under uncertainty

- Minimize the expected total buffer contents after the scheduled horizon
 – Gives maximum expected throughput
- Constrain assignments as in DEEA to satisfy delay constraints

$$\langle L \rangle = \sum_{u=1}^{U} \sum_{n_u=0}^{\infty} \sum_{x_{ut}=0}^{\infty} p(n_u|I) p(x_{ut}|I) g\left(S_u + n_u - \sum_{t=1}^{T} x_{ut}\right)$$

g(x) = x if x > 0, otherwise g(x) = 0



Traffic prediction

- We wish to determine p(n|I) given
 I = past arrival statistics
- Imagine using histograms
 - Too few observations in comparison to possible inflow sizes
- Instead, partition the inflow-axis into a number of 'bins'
 - Count arrivals within each bin
 - Adapt the bin size to obtain high resolution at intervals of high intensity and lower elsewhere



Traffic prediction – Bin probabilities



Using K bins and letting

- $-m_k = past number of arrivals of size within bin k$
- M = total number of observations

we have (after some calculations...) $p(n \in bin \ k \mid m_k, M, I) = \frac{m_k + 1}{M + K}$



Traffic prediction – Adaptation

- Based on the bin probabilities, how do we adapt the bin positions and sizes?
 - Wish to have a quantized distribution which is as close to the exact distribution as possible.
- Formally, we wish to maximize the mutual information between the two distributions

Theorem:

Maximizing the mutual information is equivalent to maximizing the entropy of the bin probability distribution.



Traffic prediction – Adaptation

Proof:

$$I(k,n) = \sum_{k=1}^{K} \sum_{n=n_{min}}^{n_{max}} p(nk) \log \frac{p(nk)}{p(n)p(k)}$$
$$= \sum_{k=1}^{K} \sum_{n=n_{min}}^{n_{max}} p(nk) \log \frac{p(k \mid n)}{p(k)}$$
$$= -\sum_{k=1}^{K} \sum_{n \in \text{bin } k} p(n \mid k)p(k) \log p(k)$$
$$= -\sum_{k=1}^{K} p(k) \log p(k)$$



Traffic prediction

- The optimum bin partition is adapted according to the M most recent arrivals:
 - Assume a uniform probability distribution within each bin,
 p(n) = bin probability / bin width
 - Redistribute the bins so that each bin has equal probability mass (=max entropy)
 - Approximate low-complexity solution requires single sweep over the possible arrival sizes.



Traffic prediction – Simulation set-up

- K = 6 bins,
- M = 100 observations between bin updates
- Min arrival rate=0, Max=100 [bit/time unit]
- Arrivals generated as
 - 50% 1-bit packets,
 - 50% 7-bit packets
 - (switching between 2 fixed rates with equal frequency)



Traffic prediction – Results





3. MIMO-Multiuser Diversity

 With a single channel, capacity-optimal schedule seds to one user at a time (Knopp&Humblet -95)

$C \sim \log(\log U)$

- Recently, the Shannon capacity for the MIMO case has been found (Caire&Shamai 2001, Vishwanath&Jindal&Goldsmith 2002, Viswanath&Tse 2002)
- The capacity-optimal scheme requires full channel knowledge at the transmitter and is achieved by Costa-precoding (extremely complex)
- Single-user transmission no longer optimal!



MIMO-Multiuser Diversity

- Consider M Tx antennas, and 1 Rx antenna and full channel knowledge at transmitter and receiver
- Then, capacity for Gaussian MIMO broadcast channel for large U is
 C ~ M log(log U)

(Sharif&Hassibi -03, submitted)

- Sharif & Hassibi proposes using M random beams where each user feeds back best beam and SINR
 - approaches capacity when $U \rightarrow \infty$



A simple MIMO-multiuser approach

- Transmit to different users on each antenna
- Let each user feed back best SINR and best antenna
- On each antenna, transmit to the user with highest SINR



A simple MIMO-multiuser approach

 Assuming M = 2 Tx antennas, and that noise is small compared to the interference from the other antenna, the rate is

$$R_{u} \leq \log \left(1 + \frac{|h_{1u}|^{2}}{\sigma^{2} + |h_{2u}|^{2}} \right)$$

$$\approx \log \left(1 + \frac{|h_{1u}|^{2}}{|h_{2u}|^{2}} \right)$$



A simple MIMO-multiuser approach

- Assume independent, identicallydistributed flat Rayleigh fading on both antennas. (Exponentially distributed |h|²).
- The distribution of a ratio of two exponentially distributed numbers with the same mean is independent of that mean.

Mathias Johansson Signals & Systems Group Thus, the performance will not depend on the average channel gain!

– Unfairness no longer a big problem



• The expected throughput depends on the probability for having high gain on one antenna, low gain on the other.

$$R_u \approx \log\left(1 + \frac{|h_{1u}|^2}{|h_{2u}|^2}\right)$$

- More users ⇔ Higher probability!
- Compare the expected throughput of this scheme to that of
 - always sending only on antenna 1,
 - sending only to the best user on the best antenna (equal to one-antenna case with twice as many users)







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Multiuser-diversity gain







Concluding remarks

- Real-time traffic and multiuser diversity is compatible for large U

 Reason is log(logU) behavior of rate
- Traffic prediction will be required if scheduling over longer time horizons
 - presented new adaptive method
- Multiuser-MIMO requires transmitting mote than one data stream at a time to approach capacity
 - simple interference-channel approach seems promising